

IMO 2012 Problems



53rd International Mathematical Olympiad
MAR DEL PLATA - ARGENTINA

Language: English

Day: 1

Tuesday, July 10, 2012

Problem 1. Given triangle ABC the point J is the centre of the excircle opposite the vertex A . This excircle is tangent to the side BC at M , and to the lines AB and AC at K and L , respectively. The lines LM and BJ meet at F , and the lines KM and CJ meet at G . Let S be the point of intersection of the lines AF and BC , and let T be the point of intersection of the lines AG and BC .

Prove that M is the midpoint of ST .

(The *excircle* of ABC opposite the vertex A is the circle that is tangent to the line segment BC , to the ray AB beyond B , and to the ray AC beyond C .)

Problem 2. Let $n \geq 3$ be an integer, and let a_2, a_3, \dots, a_n be positive real numbers such that $a_2 a_3 \cdots a_n = 1$. Prove that

$$(1 + a_2)^2 (1 + a_3)^3 \cdots (1 + a_n)^n > n^n.$$

Problem 3. The *liar's guessing game* is a game played between two players A and B . The rules of the game depend on two positive integers k and n which are known to both players.

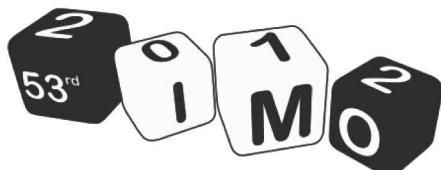
At the start of the game A chooses integers x and N with $1 \leq x \leq N$. Player A keeps x secret, and truthfully tells N to player B . Player B now tries to obtain information about x by asking player A questions as follows: each question consists of B specifying an arbitrary set S of positive integers (possibly one specified in some previous question), and asking A whether x belongs to S . Player B may ask as many such questions as he wishes. After each question, player A must immediately answer it with *yes* or *no*, but is allowed to lie as many times as she wants; the only restriction is that, among any $k+1$ consecutive answers, at least one answer must be truthful.

After B has asked as many questions as he wants, he must specify a set X of at most n positive integers. If x belongs to X , then B wins; otherwise, he loses. Prove that:

1. If $n \geq 2^k$, then B can guarantee a win.
2. For all sufficiently large k , there exists an integer $n \geq 1.99^k$ such that B cannot guarantee a win.

Language: English

Time: 4 hours and 30 minutes
Each problem is worth 7 points



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Day: 2

Wednesday, July 11, 2012

Problem 4. Find all functions $f: \mathbb{Z} \rightarrow \mathbb{Z}$ such that, for all integers a, b, c that satisfy $a+b+c=0$, the following equality holds:

$$f(a)^2 + f(b)^2 + f(c)^2 = 2f(a)f(b) + 2f(b)f(c) + 2f(c)f(a).$$

(Here \mathbb{Z} denotes the set of integers.)

Problem 5. Let ABC be a triangle with $\angle BCA = 90^\circ$, and let D be the foot of the altitude from C . Let X be a point in the interior of the segment CD . Let K be the point on the segment AX such that $BK = BC$. Similarly, let L be the point on the segment BX such that $AL = AC$. Let M be the point of intersection of AL and BK .

Show that $MK = ML$.

Problem 6. Find all positive integers n for which there exist non-negative integers a_1, a_2, \dots, a_n such that

$$\frac{1}{2^{a_1}} + \frac{1}{2^{a_2}} + \cdots + \frac{1}{2^{a_n}} = \frac{1}{3^{a_1}} + \frac{2}{3^{a_2}} + \cdots + \frac{n}{3^{a_n}} = 1.$$

Language: English

Time: 4 hours and 30 minutes
Each problem is worth 7 points



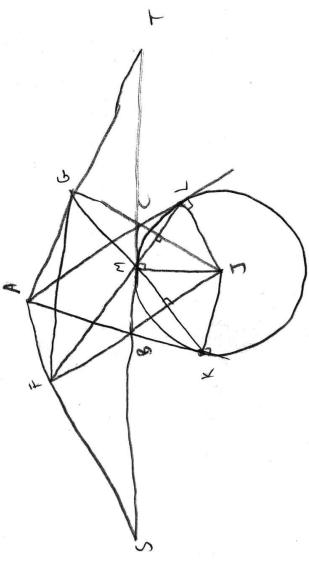
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Note that $\triangle MKJ$ and $\triangle MLJ$ are kites ($BM = BK$, $MK = CL$, $JK = JM = JL$)

$MK \perp BJ$ and $ML \perp CJ$

$\therefore M$ is the orthocenter of $\triangle FGJ \Rightarrow JM \perp FG$

But $JM \perp BC \Rightarrow BC \parallel FG$.

So now it suffices to show AM bisects FG by homothety.

$\angle FJK = \angle FJM \Rightarrow FGJK$ is cyclic

similarly $FGJL$ is cyclic.

Also note that $\angle AKJ = \angle ALJ = 90^\circ \Rightarrow AKJL$ is cyclic.

$\therefore AFKJL$ is cyclic.

$\therefore \angle AFG = \angle AKG = \angle BKM = \angle BMK = \angle FGM$

$\therefore AF \parallel GM$.

similarly $AG \parallel FM$ so $AFMG$ is a parallelogram.

$\therefore AM$ bisects $FG \Rightarrow M$ is midpoint of ST .

$$\text{Note that } 1 + q_k = \underbrace{\frac{1}{k-1} + \frac{1}{k-1} + \dots + \frac{1}{k-1}}_{k-1 \text{ of them}} + q_k \geq k \sqrt[k]{\frac{q_k}{(k-1)^{k-1}}} \quad \begin{matrix} \text{AM-GM} \\ \cancel{k-2} \\ 2 \leq k \leq n \end{matrix}$$

$$\text{So } (1 + q_k)^k \geq k^k \frac{k^k}{(k-1)^{k-1}} q_k \quad \text{with equality iff } q_k = \frac{1}{k-1}.$$

$$\begin{aligned} \therefore \prod_{i=2}^n (1 + q_i)^i &\geq \prod_{i=2}^n \frac{i^i}{(i-1)^{i-1}} q_i \\ &= \frac{2^2}{1} \cdot \frac{3^3}{2^2} \cdot \frac{4^4}{3^3} \cdots \cdot \frac{n^n}{(n-1)^{n-1}} \cdot \frac{n!}{i=2} q_i \end{aligned}$$

$$= n^n \quad \text{with equality iff } q_i = \frac{1}{i-1}.$$

~~If equality holds,~~ Equality does not hold because $\prod q_i = 1 \times \frac{1}{2} \times \frac{1}{3} \times \dots \times \frac{1}{n-1}$
 $\neq 1$ for $n \geq 2$.

$\therefore \prod_{i=2}^n (1 + q_i)^i > n^n$ which is what we want.

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3.1 We consider THE following GAME:

~~Answers~~

Let k be a positive integer known to A & B.

A thinks of a no. x with $1 \leq x \leq 2^k$.

B asks questions by the same way as the light's guessing game
(choose a set S , ask A if $x \in S$ light blinks).

But now B can only ask k questions, and among them, A must
answer truthfully at least one.

After the question and answer period, B must specify an integer
 y with $1 \leq y \leq 2^k$.

If $y \neq x$ B wins and B loses otherwise.

~~See~~

We claim that B can always guarantee a win.

We show this by induction.

$k=1$: B asks if $x \in \{1\}$.

If A replies yes, it must be true so B choose $y=2$.

" " no, it must be true so B choose $y=1$.

So B will win.

Now suppose true until some k , and we want to show for $k+1$.

Consider B asks if $x \in \{1, 2, \dots, 2^k\}$.

Suppose A replies yes.

Then B ~~protests~~ pretends A is lying and asks the next k
questions according to our induction hypothesis where $x \in \{2^{k+1}, 2^{k+2}, \dots, 2^{k+1}\}$.

So we are done.

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(we can imagine a different game where $\sum_{i=1}^k x_i \leq x \leq \sum_{i=1}^k x_i + 1$ and B must choose $x \in \{x_i + 1\}$)
it is the same game as long as x can be 2^k different possible values).

So now B has a no. y with $2^{k+1} \leq y \leq 2^{k+1}$, and $x \neq y$ confirmed, when A lies or lies

i.e. If $x \in \{1, 2, \dots, 2^k\}$, then $x \neq y$

if $x \in \{2^k + 1, 2^k + 2, \dots, 2^{k+1}\}$, then clearly $x \neq y$.

$x \neq y$ is confirmed, and B wins.

Similarly if A replies no to $x \in \{1, 2, \dots, 2^k\}$, it is the same as

A replying yes to $x \in \{2^{k+1}, \dots, 2^{k+1}\}$. same as $x \in \{1, 2, \dots, 2^k\}$ essentially

i.e. B always has a winning strategy

so by induction our claim is true for all k , that B can win.

Back to our question, we want to make B win by eliminating the no. of possible values of x from N to at most 2^k , if $N \leq 2^k$ done. Suppose $N \geq 2^{k+1}$.
If we can reduce 2^{k+1} to 2^k , then we can reduce N to 2^k by reducing
N to $N-1$ to $N-2$ to ... to 2^k .

So we just have to show when $N = 2^{k+1}$.

B asks A if ~~$x \in \{2^{k+1}\}$~~ $x \in \{2^{k+1}\}$ (at least $k+1$)

B asks A if ~~$x \in \{2^k\}$~~ many times, if A keeps answering no, then since one of the nos is $x \neq x$, $x \neq x$ and so we reduced 2^{k+1} possibilities to 2^k .

Suppose A answered yes or least once, B immediately stop asking A that.

Now we play the previous game where B asks k questions, with $1 \leq x \leq 2^k$.
then $\exists y$ with $1 \leq y \leq 2^k$ and $x \neq y$, given that $x \in \{2^{k+1}\}$ is a lie.

If it is a lie, x $\neq y$.

So $x \neq y$ confirmed so we reduced 2^{k+1} possibilities to 2^k .

So we are done.



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Mathematical Medley

□

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3.2 Choose $N = n+1$. Let d be a real no. s.t. $1.99 < d < 2$.

When A says $x \in S$ or $x \notin S$, he claims that

$x \neq y \forall y \in S$ or $x \neq y \forall y \in S$.

Right after the n^{th} question, let $v_{n,m}$ be the largest no. P satisfying $1 \leq q \leq N$. Right after the n^{th} question, let s_m be the largest no. P satisfying

for the last P questions, A claims that $x \in S$.

Then $v_{n,p} = 0 \quad \forall n \in \mathbb{N}, 1 \leq q \leq N$.

Let $v_{n,m} = d v_{n,m}$, and $s_m = \sum_{n=1}^m v_{n,m}$.

then $s_0 = N$.

Now A's strategy is to consider yes or no, whichever leads to the lowest possible s_m score.

Note that $v_{n,m+1} = \begin{cases} v_{n,m} + 1 & \text{if A claims } x \in S \text{ on the } m^{\text{th}} \text{ turn.} \\ 0 & \text{otherwise.} \end{cases}$

$$v_{n,m+1} = \begin{cases} d v_{n,m} & \text{if } \dots \\ 0 & \text{if } \dots \end{cases}$$

Suppose at the m^{th} turn A says yes and s_m will have a value s'_m .

||| no and s_m will have a value s''_m .

$$\text{Then } s'_m + s''_m = d s_{m-1} + N.$$

so the lowest possible s_m is $\leq \frac{d}{2} s_{m-1} + \frac{N}{2}$.

$$\Rightarrow s_m \leq \frac{d}{2} s_{m-1} + \frac{N}{2} \quad \boxed{s_m - \frac{dN}{2-d} \leq \frac{d}{2} (s_{m-1} - \frac{dN}{2-d})}$$

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Note that $s_0 - \frac{dN}{2-d} < 0 \quad \text{so} \quad s_m - \frac{dN}{2-d} < 0 \quad \forall m$.

$$\Rightarrow s_m < \frac{dN}{2-d} \quad \forall m$$

Suppose there is a no. k where A claims $x \in S$ in k consecutive turns.

Then there is a no. m where $s_m \geq d^{k+1}$.

$$\text{So } d^{k+1} < \frac{dN}{2-d}.$$

We can choose $1 + 1.99^k < N < 1.999^k$ for sufficiently large k .

$$\text{Then } d^{k+1} < \frac{d}{2-d} \times 1.999^k \Rightarrow \left(\frac{d}{1.999}\right)^{k+1} < \frac{d}{2-d}.$$

But when we choose $d = 1.999907026031062$, the above is not true

for sufficiently large k .

\therefore A sufficiently large k , A claims that $x \in S$ in at most k consecutive turns.

A can choose $d = 1.999907026031062$, so x can be any value from 1 to N and it will be truthful among any $k+1$ consecutive answers. So B cannot guarantee a win.

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Set $a+b+c=0$, we get $f(0)^2=0 \Rightarrow f(0)=0$.

$$\begin{aligned} \text{Set } c=0, b=-a, \text{ we get } & f(a)+f(-a)-2f(a)f(-a)=0 \\ & \Rightarrow (f(a)-f(-a))^2=0 \\ & \Rightarrow f(a)=f(-a). \end{aligned}$$

Set $c=-a-b$, then since $f(a+b)=f(-a-b)$,

$$\begin{aligned} f(a)^2 + f(b)^2 + f(a+b)^2 &= 2f(a)f(b) + 2f(b)f(c) + 2f(c)f(a) \\ \Leftrightarrow (f(a+b) - f(a) - f(b))^2 &= 4f(a)f(b). \end{aligned}$$

This means that $\forall a, b \in \mathbb{Z}$, $f(a)f(b) \geq 0$. So f is always non-negative or non-positive.

WLOG assume $f(x) > 0 \forall x$, since if $f(x) \leq 0 \forall x$, replace $f(x)$ by $-f(x)$.

$$\begin{aligned} \therefore f(a+b) - f(a) - f(b) &= \pm 2\sqrt{f(a)f(b)} \\ \Leftrightarrow f(a+b) &= (\sqrt{f(a)} \pm \sqrt{f(b)})^2 \\ \Leftrightarrow \sqrt{f(a+b)} &= \pm \sqrt{f(a)} \pm \sqrt{f(b)} \quad (\text{with } \pm \text{ may not be the same sign}). \end{aligned}$$

Let $g: \mathbb{Z} \rightarrow \mathbb{R}$ where $g(x) = \sqrt{f(x)} \forall x$.

Now the original condition is equivalent to $g(a+b) \leq g(a) + g(b) \quad \forall a, b$ and $g(0)^2 \in \mathbb{Z}$
and $g(0) = g(-a)$ and $g(a) \geq 0$.

So we only need to check these 4 conditions in order to satisfy the original condition.

But just to be safe let's just throw in $g(0)=0$
to have 5 conditions.

Suppose $\exists z \neq 0$ such that $g(z) = 0$ for some $z \neq 0$.

Then $\exists x, y \in \mathbb{Z}$ such that $g(x+z) = \pm g(x) = \pm g(y) = g(x) = g(y) \geq 0$. So g has period 2.

Setting $x=z, 2z, 3z, \dots$ gives $g(z) = g(2z) = g(3z) = \dots = 0$.

$\therefore g(nz) = 0 \quad \forall n \in \mathbb{Z}$ (using $g(m) = g(m-n)$).

Let $g(1) = k$. If $k=0$ then $g(x)=0 \forall x$ which works.

Otherwise $k > 0$, then $k \in \mathbb{Z}$.

Then $g(x+1) = \pm g(x) = k$. ~~so we can easily induction stand~~

Suppose $g(1)=k$, $g(2)=\pm k, \dots, g(a)=ak$ for some $a \geq 1$,

and $g(a+1)=ak-k$.

It is not hard to see α exists from $g(a+\alpha) = \pm g(a) = k$ or $g(x) = |x|k \forall x$.

Easy to check $g(x) = |x|k$ works, for $\frac{x}{2} \in \mathbb{Z}$.

$$\begin{aligned} \text{If } a \geq 4, \quad g(a+2) &= \pm g(a+1) \pm k = ak-2k \text{ or } ak \\ g(a+1) = ak-k = \pm 2k \pm (a-1)k &= \pm g(a) \pm 2k = ak-2k \text{ or } ak+2k \\ \text{impossible}. \end{aligned}$$

$\therefore g(a+2) = \pm 2k$ ~~not possible~~.

~~g(a+2) = 2k~~

Case 1: $a=1$

$g(1)=k$, $g(2)=0$, so $g(x)=g(x+2)$ is periodic.

So $g(x) = \begin{cases} k & \text{when } x \text{ is odd} \\ 0 & \text{when } x \text{ is even} \end{cases}$ easy to check it works.

Case 2: $a=2$

$g(1)=k$, $g(2)=2k$, $g(3)=k$, $g(4)=2k$ $\therefore g(a)=0$.

So $g(x) = \begin{cases} 0 & x \leq 0 \\ k & x=1, 3 \pmod{4} \\ 2k & x \geq 2 \end{cases}$

To check it works, ~~constant~~ $g(6+4) = \pm g(2) \neq g(6)$ only for $a, b \in \{0, 1, 2, 3\}$, which is easy.

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Consider $N \in \min\{q_1, q_2, \dots, q_n\}$ then
 $3^{N-q_1} + 2 \times 3^{N-q_2} + \dots + n \times 3^{N-q_n} - 3^N \equiv 1 \pmod{4}$.

$$\therefore \frac{n(n+1)}{2} \equiv 1 \pmod{4} \Rightarrow n \equiv 1, 2 \pmod{4}.$$

Consider small values of n , easy to check these are solutions:

$\{q_n\}$ in that order	
0	0
1	1, 1
2	2, 2, 2, 3, 3
5	2, 3, 3, 3, 4, 4, 4, 4, 4
6	2, 2, 2, 3, 3, 3, 3, 3, 3
9	2, 3, 3, 3, 4, 4, 4, 4, 4
10	2, 3, 3, 4, 4, 4, 4, 4, 4

(trust me)

We claim that all n of the form $1, 2 \pmod{4}$ works. We show this by extending the construction for n to $n+8$.

$$\text{Note that } \frac{1}{2^x} = \frac{1}{2^{x+2}} + \frac{1}{2^{x+3}} + \frac{1}{2^{x+4}} * 6.$$

Case 1: $n \equiv 0 \pmod{4}$

$$\begin{aligned} \text{Consider } \frac{k}{3^a} &= \frac{k}{3^{x+2}} + \frac{n+1}{3^{x+3}} + \frac{n+2}{3^{x+4}} + \frac{n+3}{3^{x+5}} + \frac{n+4}{3^{x+6}} + \\ &\quad + \frac{n+2}{3^{x+3}} + \frac{n+4}{3^{x+5}} \\ \Leftrightarrow 4k &= n+4. \end{aligned}$$

So k exists (in integer) and we just extended n to $n+8$.

Because if (b_1, b_2, \dots, b_n) is a solution for n , $(b_1, b_2, \dots, b_{n-1}, b_{n+2}, b_{n+3}, \dots, b_n, b_{n+4}, b_{n+5}, b_{n+6}, b_{n+7}, b_{n+8})$ is a solution for $n+8$.

is a solution for $n+8$)

case 2: $n \equiv 1 \pmod{4}$

$$\text{Consider } \frac{k}{3^a} = \frac{k}{3^{x+2}} + \frac{n+3+n+4+n+5+n+6+n+7+n+8}{3^{x+4}} + \frac{n+1}{3^{x+3}} + \frac{n+2}{3^{x+2}}$$

$$\Leftrightarrow 4k = n+3 \quad \text{so } k \text{ exists, and similarly we extended } n \text{ to } n+8.$$

case 3: $n \equiv 2 \pmod{4}$

$$\text{Consider } \frac{k}{3^a} = \frac{k}{3^{x+2}} + \frac{n+1+n+2+n+3+n+5+n+6+n+7}{3^{x+4}} + \frac{n+4}{3^{x+3}} + \frac{n+3}{3^{x+2}}$$

$$\Leftrightarrow 4k = n+6 \quad \text{so } k \text{ exists, and similarly done.}$$

case 4: $n \equiv 3 \pmod{4}$.

I just realized we don't need case 4 & 1, since n can only be $1, 2 \pmod{4}$.

all possible n is $\frac{1}{2}$ (mod 4).

Using this method, we can extend $n=1$ to $n=9$

$$0 \rightarrow 2, 3, 2, 4, 4, 4, 4, 4, 4$$

extend $n=2$ to 10

$$1, 1 \rightarrow 1, 3, 5, 5, 4, 5, 5, 5, 3$$

$$1, 3, 5, 5, 4, 5, 5, 3 \rightarrow 1, 3, 5, 7, 5, 4, 5, 5, 3, 9, 9, 8, 9, 9, 7$$

because if (b_1, b_2, \dots, b_n) is a solution for n , $(b_1, b_2, \dots, b_{n-1}, b_{n+2}, b_{n+3}, \dots, b_n, b_{n+4}, b_{n+5}, b_{n+6}, b_{n+7}, b_{n+8})$ is a solution for $n+8$.